

## Section 7: A lower bound algorithm

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As previously, we consider a block structure

$$\Delta = \{ \text{diag}(\delta_1 I_{r_1}, \dots, \delta_n I_{r_n}, \Delta_{S+1}, \dots, \Delta_{S+F}) : \delta_j \in \mathbb{C}, \Delta_{S+j} \in \mathbb{C}^{m_j \times m_j} \},$$

where

$$r_1 + \dots + r_S + m_1 + \dots + m_F = n.$$

Recall that for  $M \in \mathbb{C}^{n \times n}$

$$\mu_{\Delta}(M) = \max_{\Delta \in \mathbf{B}_{\Delta}} \rho(\Delta M) = \max_{Q \in \mathbf{Q}} \rho(QM)$$

and

$$\mathbf{B}_{\Delta} = \{ \Delta \in \Delta : \bar{\sigma}(\Delta) \leq 1 \},$$

$$\mathbf{Q} = \{ Q \in \Delta : Q \text{ is unitary} \},$$

$$\mathbf{D} = \{ \text{diag}(D_1, \dots, D_S, d_{S+1} I_{m_1}, \dots, d_{S+F} I_{m_F}) : \\ D_j \in \mathbb{C}^{r_j \times r_j}, D_j > 0, d_{S+j} \in \mathbb{R}_{>0} \}.$$

Note that every  $D \in \mathbf{D}$  and  $\Delta \in \Delta$  commute.

## Observation

Let  $M \in \mathbb{C}^{n \times n}$ ,  $\beta > 0$ . If there exist  $Q \in \mathbf{Q}$ ,  $D \in \mathbf{D}$ ,  $\xi \in \mathbb{C}^n$ ,  $\|\xi\| = 1$  such that

$$QD^{\frac{1}{2}}MD^{-\frac{1}{2}}\xi = \beta\xi, \quad (1)$$

$$D^{-\frac{1}{2}}M^*D^{\frac{1}{2}}Q^*\xi = \beta\xi, \quad (2)$$

then

$$\beta \leq \mu_{\Delta}(M).$$

Putting  $\eta = \|D^{-\frac{1}{2}}\xi\|^{-1}D^{-\frac{1}{2}}\xi$  we obtain that (1) and (2) are equivalent to

$$\exists \eta \in \mathbb{C}^n, \|\eta\| = 1 : QM\eta = \beta\eta, \quad (QM)^*D\eta = \beta D\eta. \quad (3)$$

**Proof of Observation:**  $\beta$  is an eigenvalue of  $QM$ , so  $\beta \leq \rho(QM) \leq \mu_{\Delta}(M)$  (actually, here we use only (1)).

## Example

$S = 1$ ,  $F = 0$ ,  $r_1 = 2$ ,  $\Delta = \{\text{diag}(\delta, \delta) : \delta \in \mathbb{C}\}$ .

$$M := \begin{bmatrix} 3 & -1 \\ 1 & 1 \end{bmatrix}.$$

We have  $\det(qM - \lambda) = (2q - \lambda)^2$  and hence

$$\mu_{\Delta}(M) = \max_{|q|=1} \rho(QM) = 2,$$

but there is no  $\beta > 0$  which satisfies (1) and (2) (although there is much  $\beta$ 's satisfying (1) with some  $Q$ ,  $D$ ,  $\xi$ ).

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Indeed, (1) and (2) hold iff (3) hold with some  $Q \in \mathbf{Q}$ ,  $D \in \mathbf{D}$ ,  $\|\eta\| = 1$ . This gives  $Q = I_2$ ,  $\beta = 2$  and

$$\eta \in \mathbb{R} \cdot (1, 1), \quad D\eta \in \mathbb{R} \cdot (1, -1).$$

Hence  $0 = \langle D\eta, \eta \rangle > 0$ , a contradiction.

### Example

$S = 1$ ,  $F = 0$ ,  $r_1 = 3$ ,  $\Delta = \{\text{diag}(\delta, \delta, \delta) : \delta \in \mathbb{C}\}$ .

$$M := \begin{bmatrix} 3 & -1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}.$$

We have  $\det(qM - \lambda) = (2q - \lambda)^2(q - \lambda)$  and hence

$$\mu_{\Delta}(M) = \max_{|q|=1} \rho(QM) = 2,$$

but the only  $\beta > 0$  which satisfies (1) and (2) is  $\beta = 1$ .

Let  $x, y \in \mathbb{C}^n$ . Write them compatibly to  $\Delta$ :

$$x = (x_1, \dots, x_S, x_{S+1}, \dots, x_{S+F}) \in \mathbb{C}^{r_1} \times \dots \times \mathbb{C}^{r_S} \times \mathbb{C}^{m_1} \times \dots \times \mathbb{C}^{m_F} = \mathbb{C}^n,$$

$$y = (y_1, \dots, y_S, y_{S+1}, \dots, y_{S+F}) \in \mathbb{C}^{r_1} \times \dots \times \mathbb{C}^{r_S} \times \mathbb{C}^{m_1} \times \dots \times \mathbb{C}^{m_F} = \mathbb{C}^n.$$

We say that the pair  $(x, y)$  is *non-degenerate* (with respect to  $\Delta$ ), if

$$\langle x_1, y_1 \rangle, \dots, \langle x_S, y_S \rangle \neq 0 \text{ and } x_{S+1}, y_{S+1} \neq 0, \dots, x_{S+F}, y_{S+F} \neq 0.$$

### Theorem

Let  $M \in \mathbb{C}^{n \times n}$ . Assume that there exists  $Q_0 \in \mathbf{Q}$  such that:

- $\rho(Q_0 M) = \max_{Q \in \mathbf{Q}} \rho(QM) > 0$ ,
- $\rho(Q_0 M)$  is a distinct eigenvalue of  $Q_0 M$ ,
- there exist a non-degenerate pair  $(x, y) \in \mathbb{C}^n \times \mathbb{C}^n$  such that

$$Q_0 M x = \rho(Q_0 M) x, \quad (Q_0 M)^* y = \rho(Q_0 M) y.$$

Then there exists  $D \in \mathbf{D}$  and  $\xi \in \mathbb{C}^n$ ,  $\|\xi\| = 1$  such that

$$\begin{aligned} Q_0 D^{\frac{1}{2}} M D^{-\frac{1}{2}} \xi &= \mu_{\Delta}(M) \xi, \\ D^{-\frac{1}{2}} M^* D^{\frac{1}{2}} Q_0^* \xi &= \mu_{\Delta}(M) \xi. \end{aligned}$$

## Lemma

Let  $x, y \in \mathbb{C}^n$  be non-degenerate vectors. Then the following are equivalent:

- 1 there exists  $D \in \mathbf{D}$  such that  $y = Dx$ .
- 2 for every  $G \in \Delta$  such that  $G + G^* \leq 0$ ,  $GG^* = G^*G$  there is  $\operatorname{Re} \langle Gx, y \rangle \leq 0$ .

**Proof of the lemma.** (1)  $\Rightarrow$  (2):  $D$  and  $G$  commute and  $y = Dx$ , so

$$\langle Gx, y \rangle = \langle Gx, D^{\frac{1}{2}} D^{\frac{1}{2}} x \rangle = \langle D^{\frac{1}{2}} Gx, D^{\frac{1}{2}} x \rangle = \langle GD^{\frac{1}{2}} x, D^{\frac{1}{2}} x \rangle,$$

$$\overline{\langle Gx, y \rangle} = \langle D^{\frac{1}{2}} x, GD^{\frac{1}{2}} x \rangle = \langle G^* D^{\frac{1}{2}} x, D^{\frac{1}{2}} x \rangle.$$

Hence

$$2\operatorname{Re} \langle Gx, y \rangle = \langle (G + G^*) D^{\frac{1}{2}} x, D^{\frac{1}{2}} x \rangle \leq 0.$$

(2)  $\Rightarrow$  (1): It suffices to prove it for each block separately, so in fact we need to consider only two cases:

- $\Delta = \mathbb{C}^{n \times n}$ ,
- $\Delta = \{\operatorname{diag}(\delta, \dots, \delta) : \delta \in \mathbb{C}\}$ .



The case  $\Delta = \mathbb{C}^{n \times n}$ . Since

$$\mathbf{D} = \{\text{diag}(d, \dots, d) : d \in \mathbb{R}_{>0}\},$$

we have to prove that  $y = dx$  for some  $d > 0$ . We may assume that

$$\|x\| = \|y\| = 1.$$

There exists an unitary matrix  $U \in \mathbb{C}^{n \times n}$  such that

$$Ux = y.$$

We have

$$U = P^*JP$$

for some unitary matrices  $P, J \in \mathbb{C}^{n \times n}$ ,  $J = \text{diag}(\zeta_1, \dots, \zeta_n)$ ,  $\zeta_j \in \mathbb{T}$ . Write

$$Px = (a_1, \dots, a_n).$$

There is

$$Py = JPx = (\zeta_1 a_1, \dots, \zeta_n a_n).$$

It suffices to show that  $a_j \zeta_j = a_j$  for every  $j$ .

Without loss of generality:  $j = 1$ .

Suppose that  $a_1\zeta_1 \neq a_1$ , i.e.  $a_1 \neq 0$  and  $\zeta_1 \neq 1$ . There exists  $\eta_1 \in \mathbb{T}$  such that

$$\operatorname{Re} \eta_1 \leq 0, \operatorname{Re}(\eta_1 \bar{\zeta}_1) > 0.$$

Let

$$G := P^* \operatorname{diag}(\eta_1, 0, \dots, 0) P$$

We have

$$\begin{aligned} 0 \geq \operatorname{Re} \langle Gx, y \rangle &= \langle \operatorname{diag}(\eta_1, 0, \dots, 0) Px, Py \rangle \\ &= \langle (\eta_1 a_1, 0, \dots, 0), (a_1 \zeta_1, \dots, a_n \zeta_n) \rangle \\ &= |a_1|^2 \operatorname{Re}(\eta_1 \bar{\zeta}_1) > 0, \end{aligned}$$

a contradiction.

**The case**  $\Delta = \{\operatorname{diag}(\delta, \dots, \delta) : \delta \in \mathbb{C}\}$ . We have

$$\mathbf{D} = \{D \in \mathbb{C}^{n \times n} : D > 0\},$$

Existence of  $D$  such that  $Dx = y$  follows from the fact that  $\langle x, y \rangle > 0$ .

**Proof of the theorem. Step 1:** If  $Q_0 = I$ , then there exists  $D \in \mathbf{D}$  such that

$$y = Dx.$$

**Proof of step 1:** Let  $G \in \Delta$ ,  $G + G^* \leq 0$ ,  $GG^* = G^*G$ . It suffices to show that

$$\operatorname{Re} \langle Gx, y \rangle \leq 0.$$

Define

$$W(t) := e^{tG} M, \quad t \geq 0.$$

$\rho(M)$  is a distinct eigenvalue of  $M$ , so  $\langle x, y \rangle \neq 0$  and we may assume  $\langle x, y \rangle = 1$ .

There are  $\epsilon > 0$  and  $\mathbb{R}$ -analytic maps  $X, Y : (-\epsilon, \epsilon) \rightarrow \mathbb{C}^n$ ,  $\lambda : (-\epsilon, \epsilon) \rightarrow \mathbb{C}$  s.t.

- $\lambda(0) = \rho(W(0)) = \rho(M)$ ,  $X(0) = x$ ,  $Y(0) = y$ ,
- $\lambda(t)$  is a distinct eigenvalue of  $W(t)$ ,
- $|\lambda(t)| = \rho(W(t))$ ,
- $W(t)X(t) = \lambda(t)X(t)$ ,  $W(t)^*Y(t) = \overline{\lambda(t)}Y(t)$ ,
- $\langle X(t), Y(t) \rangle = 1$ .

Since  $GG^* = G^*G$ , we have

$$e^{tG} \in \mathbf{B}_\Delta, \quad t \geq 0,$$

so the function  $[0, \epsilon) \ni t \mapsto \operatorname{Re} \lambda(t)$  has a maximum at  $t = 0$ . There holds

$$\lambda'(0) = \langle W'(0)x, y \rangle.$$

Hence

$$0 \geq \operatorname{Re} \lambda'(0) = \operatorname{Re} \langle W'(0)x, y \rangle = \operatorname{Re} \langle GMx, y \rangle = \rho(M) \operatorname{Re} \langle Gx, y \rangle.$$

**Step 2:** We prove the conclusion.

By **step 1** for  $\tilde{M} := Q_0M$  we have  $y = Dx$  for some  $D \in \mathbf{D}$ . Put

$$\xi = D^{\frac{1}{2}}x = D^{-\frac{1}{2}}y.$$

As  $Q_0$  and  $D^{\frac{1}{2}}$  commute, we have

$$Q_0 D^{\frac{1}{2}} M D^{-\frac{1}{2}} \xi = D^{\frac{1}{2}} Q_0 M D^{-\frac{1}{2}} D^{\frac{1}{2}} x = \rho(Q_0 M) D^{\frac{1}{2}} x = \mu_\Delta(M) \xi$$

and similarly we obtain the condition (2).

From now, for simplicity:  $S = F = 1$ ,  $r + m = n$ , i.e.

$$\begin{aligned}\Delta &= \left\{ \text{diag}(\delta_1 I_r, \Delta_2) : \delta_1 \in \mathbb{C}, \Delta_2 \in \mathbb{C}^{m \times m} \right\}, \\ \mathbf{D} &= \left\{ \text{diag}(D_1, d_2 I_m) : D_1 \in \mathbb{C}^{r \times r}, D_1 > 0, d_2 \in \mathbb{R}_{>0} \right\}.\end{aligned}$$

### Proposition

Let  $M \in \mathbb{C}^{n \times n}$ ,  $\beta > 0$ . Then there exist  $Q \in \mathbf{Q}$ ,  $D \in \mathbf{D}$  and  $\xi = (\xi_1, \xi_2) \in \mathbb{C}^n$ ,  $\|\xi\| = 1$  such that

$$\begin{aligned}QD^{\frac{1}{2}}MD^{-\frac{1}{2}}\xi &= \beta\xi, \\ D^{-\frac{1}{2}}M^*D^{\frac{1}{2}}Q^*\xi &= \beta\xi,\end{aligned}$$

iff there exist  $a = (a_1, a_2)$ ,  $w = (w_1, w_2)$ ,  $b, z \in \mathbb{C}^r \times \mathbb{C}^m$  such that  $(a, w)$  is non-degenerate and

$$\begin{aligned}\beta a &= Mb, \quad z = \left( \frac{\langle a_1, w_1 \rangle}{|\langle a_1, w_1 \rangle|} w_1, \frac{\|w_2\|}{\|a_2\|} a_2 \right), \\ \beta w &= M^*z, \quad b = \left( \frac{\langle w_1, a_1 \rangle}{|\langle w_1, a_1 \rangle|} a_1, \frac{\|a_2\|}{\|w_2\|} w_2 \right).\end{aligned}$$

**Proof of the proposition.** ( $\Rightarrow$ ) We have

$$M \left( D^{-\frac{1}{2}} \xi \right) = \beta \left( D^{-\frac{1}{2}} Q^* \xi \right), \quad M^* \left( D^{\frac{1}{2}} Q^* \xi \right) = \beta \left( D^{\frac{1}{2}} \xi \right).$$

Set

$$b := D^{-\frac{1}{2}} \xi, \quad a := D^{-\frac{1}{2}} Q^* \xi, \quad z := D^{\frac{1}{2}} Q^* \xi, \quad w := D^{\frac{1}{2}} \xi.$$

( $\Leftarrow$ ) Set

$$Q = \text{diag} \left( \frac{\langle w_1, a_1 \rangle}{|\langle w_1, a_1 \rangle|} I_r, Q_2 \right),$$

where  $Q_2 \in \mathbb{C}^{m \times m}$  is unitary and such that  $Q_2 z_2 = w_2$ ,

$$D = \text{diag} \left( D_1, \frac{\|w_2\|}{\|a_2\|} I_m \right),$$

where  $D_1 \in \mathbb{C}^{r \times r}$  is positive and such that  $D_1 a_1 = z_1$  (it exists, because  $\langle a_1, z_1 \rangle > 0$ ), and

$$\xi := D^{\frac{1}{2}} b.$$

## An iterative algorithm for a lower bound for $\mu_{\Delta}(M)$ :

Take  $a^1 = (a_1^1, a_2^1)$ ,  $w^1 = (w_1^1, w_2^1)$ ,  $b^1, z^1 \in (\mathbb{C}^r)_* \times (\mathbb{C}^m)_*$  such that

$$\langle a_1^1, w_1^1 \rangle \neq 0, \quad \|a^1\| = \|b^1\| = \|z^1\| = \|w^1\| = 1.$$

Define (assuming that the following definitions are proper)  $\tilde{\beta}_k, \hat{\beta}_k \in \mathbb{R}_{>0}$ ,  $a^k = (a_1^k, a_2^k)$ ,  $w^k = (w_1^k, w_2^k)$ ,  $b^k, z^k \in (\mathbb{C}^r)_* \times (\mathbb{C}^m)_*$  as

$$\begin{aligned}\tilde{\beta}_{k+1} a^{k+1} &:= Mb^k \text{ with } \|a^{k+1}\| = 1 \\ z^{k+1} &:= \left( \frac{\langle a_1^{k+1}, w_1^k \rangle}{|\langle a_1^{k+1}, w_1^k \rangle|} w_1^k, \frac{\|w_2^k\|}{\|a_2^{k+1}\|} a_2^{k+1} \right), \\ \hat{\beta}_{k+1} w^{k+1} &:= M^* z^{k+1} \text{ with } \|w^{k+1}\| = 1 \\ b^{k+1} &:= \left( \frac{\langle w_1^{k+1}, a_1^{k+1} \rangle}{|\langle w_1^{k+1}, a_1^{k+1} \rangle|} a_1^{k+1}, \frac{\|a_2^{k+1}\|}{\|w_2^{k+1}\|} w_2^{k+1} \right).\end{aligned}$$

We have

$$\|a^k\| = \|b^k\| = \|w^k\| = \|z^k\| = 1.$$

If  $\widehat{\beta}_k, \widetilde{\beta}_k, a_k, b_k, w_k, z_k$  converge (respectively) to some  $\widehat{\beta}, \widetilde{\beta} \in \mathbb{R}_{>0}$ ,  
 $a = (a_1, a_2), w = (w_1, w_2), b, z \in \mathbb{C}^r \times \mathbb{C}^m$  with  $(a, w)$  being non-degenerate, then

$$\widehat{\beta} = \widetilde{\beta} \leq \mu_{\Delta}(M).$$

**Proof:** We have

$$\begin{aligned} \widetilde{\beta}a &= Mb, & z &= \left( \frac{\langle a_1, w_1 \rangle}{|\langle a_1, w_1 \rangle|} w_1, \frac{\|w_2\|}{\|a_2\|} a_2 \right), \\ \widehat{\beta}w &= M^*z, & b &= \left( \frac{\langle w_1, a_1 \rangle}{|\langle w_1, a_1 \rangle|} a_1, \frac{\|a_2\|}{\|w_2\|} w_2 \right). \end{aligned}$$

Thus

$$\widetilde{\beta}\langle a, z \rangle = \langle Mb, z \rangle = \langle b, M^*z \rangle = \widehat{\beta}\langle b, w \rangle = \widehat{\beta}\langle a, z \rangle,$$

so  $\widehat{\beta} = \widetilde{\beta}$ . Now apply the previous proposition.



In practice:

- If the algorithm seems to converge, i.e. if  $\widehat{\beta}_k \approx \widetilde{\beta}_k$ , then  $\widehat{\beta}_k$  and  $\widetilde{\beta}_k$  approximate some lower bound of  $\mu_{\Delta}(M)$ .

In practice:

- If the algorithm seems to converge, i.e. if  $\hat{\beta}_k \approx \tilde{\beta}_k$ , then  $\hat{\beta}_k$  and  $\tilde{\beta}_k$  approximate some lower bound of  $\mu_{\Delta}(M)$ .
- In the opposite case we can just restart the algorithm with different initial points.

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- If the algorithm seems to converge, i.e. if  $\widehat{\beta}_k \approx \widetilde{\beta}_k$ , then  $\widehat{\beta}_k$  and  $\widetilde{\beta}_k$  approximate some lower bound of  $\mu_{\Delta}(M)$ .
- In the opposite case we can just restart the algorithm with different initial points.
- After some modifications of the algorithm it is possible to avoid the problems with definition of the sequences or with non-convergence.

Thank you for attention.